



## Ceylon at Jfokus

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**We're fans of Java**

# Disclaimer

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If I sound critical of Java (or any other language) in this presentation  
it's merely to identify problems that require solutions

Indeed, a lot of criticism of Java is IMO deeply misplaced—but that  
doesn't mean there's nothing wrong with Java!

# What is it?

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## A programming language:

- That runs on *virtual machines*
- To be specific, the Java VM, and JavaScript VMs
- Defined by a *specification*
- With a syntax that looks conventional but is actually very flexible
- With an extremely powerful and elegant type system
- With built-in *modularity*
- With its own language module and SDK
- And *excellent tooling*

# What is it?

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## Where it runs:

- On Java SE, with the Ceylon module runtime
- In any OSGi container: Eclipse, Apache Felix, WildFly, GlassFish, ...
- On Vert.x
- On Node.js
- In a web browser, with Common JS Modules (require.js)
- In a Java servlet engine, via `ceylon war` (in the next release)

# What about interop?

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## Interoperable with native code

- It can be used to build a cross-platform module that executes in both virtual machine environments, depending only on other cross-platform modules written in pure Ceylon
- Or, it can be used to write a module that targets only one of the two virtual machines and interoperates with native Java or JavaScript code for that platform
- Interoperation with JavaScript is via dynamic typing, or via writing an interface in Ceylon that ascribes static types to the JS API

# A few unique things

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## Some unique things about Ceylon

- *Designed for multiplatform use*—the language and language module completely abstract the details of the virtual machine
- *Reified generics*, along with a *typesafe metamodel* that provides access to generic type arguments at runtime
- *Union and intersection types*—the foundation for unambiguous *type inference* and *flow-sensitive typing*
- Representation and abstraction of *function and tuple types* within the type system—without an explosion of single-method interface types or `Function1`, `Function2`, `Function3`, ...
- A simple, unified type system, with elegant *syntax sugar* that helps reduce verbosity without harming readability

# Idiom #1

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## Idiom: functions with multiple outcomes

For example, an operation might return a `File`, a `Url`, or nothing:

```
//Java
Object parsePath(String path)
    throws SyntaxException { ... }
```

We can handle the different outcomes using `instanceof`, type casts, and `catch`:

```
try {
    Object result = parsePath(path);
    if (result instanceof File) {
        File file = (File) result;
        return lines(file);
    }
    if (result instanceof Url) {
        Url url = (Url) result;
        return new Request(url).execute().getContent().getLines();
    }
}
catch (SyntaxException se) { return emptyList(); }
```

# Idiom #1

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## Idiom: functions with multiple outcomes

A function with more than one “outcome” can be defined using a union type.

```
File|Path|SyntaxError parsePath(String path) => ... ;
```

We can handle the various outcomes using `switch`:

```
value result = parsePath(name);
switch (result)
case (is File) {
    return lines(result);
}
case (is Url) {
    return Request(result).execute().content.lines;
}
case (is SyntaxError) {
    return {};
}
}
```

# Idiom #1

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## Idiom: functions with multiple outcomes

We can aggregate cases using union:

```
value result = parsePath(name);
switch (result)
case (is File|Url) {
    ...
}
else {
    ...
}
```

Or, alternatively, using `if` instead of `switch`:

```
if (is File|Url result
    = parsePath(name)) {
    ...
}
```

# Idiom #2

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## Idiom: functions returning null

Example: retrieve an item from a map.

(Nothing more than a special case of multiple outcomes!)

```
Item? get(Key key) => ... ;
```

Here `Item?` literally means `Null | Item`.

```
value map = HashMap { "CET"->cst, "GMT"->gmt, "PST"->pst };
```

```
Timezone tz = map[id]; //not well-typed!
```

```
value offset = map[id].rawOffset; //not well-typed!
```

```
Timezone? tz = map[id];
```

```
value offset = (map[id] else gmt).rawOffset;
```

For a union type of this very common form, we have special syntax sugar.

# Idiom #3

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## Idiom: heterogeneous collections

What is the type of a list containing `Integers` and `Floats`?

```
//Java  
List<Number> list = Arrays.asList(1, 2, 1.0, 0.0);
```

The element type is ambiguous, so I must be explicit.

Even then I lose some information.

```
Number element = list.get(index);  
//handle which the subtypes of Number?  
//don't forget that an out of bounds  
//index results in an exception
```

# Idiom #3

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## Idiom: heterogeneous collections

With union and intersection, type inference is unambiguous!

```
value list = ArrayList { 1, 2, 1.0, 0.0 };
```

The inferred element type is `Integer|Float`, resulting in the inferred type `ArrayList<Integer|Float>`, which is a subtype of any type to which the `ArrayList` may be soundly assigned.

No loss of precision!

```
Integer|Float|Null element = list[index];  
//now I know exactly which cases I have to handle
```

# Idiom #4

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## Idiom: unions and streams

Example: the `follow()` method of `Iterable` adds an element to the start of a stream.

```
{Element|Other+} follow<Other>(Other element)
=> { element, *this };
```

The syntax `{T*}` and `{T+}` is sugar for the interface `Iterable`.

Exactly the right type pops out automatically.

```
{String*} words = { "hello", "world" };
{String?+} strings = words.follow(null);
```

(Even though I'm explicitly writing in the types, I could have let them be inferred.)

# Idiom #5

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## Idiom: intersections and streams

Example: the `coalesce()` function eliminates `null` elements from a stream.

```
{Element&Object*} coalesce<Element>({Element*} elements)  
=> { for (e in elements) if (exists e) e };
```

Again, exactly the right type pops out automatically.

```
{String?*} words = { "hello", null, "world" };  
{String*} strings = coalesce(words);
```

(Again, I could have let the types be inferred.)

# Idiom #6

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## Idiom: empty vs nonempty

Problem: the `max()` function can return `null`, but only in the case that the stream might be empty. So let's try this:

```
shared Value? max<Value>({Value*} values)
  given Value satisfies Comparable<Value> { ... }
```

What if we know it's nonempty at compile time? Do we need a separate function?

```
shared Value maxNonempty<Value>({Value+} values)
  given Value satisfies Comparable<Value> { ... }
```

Terrible! This doesn't let us abstract.

# Idiom #6

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## Idiom: empty vs nonempty

Solution: the `Iterable` type has an extra type parameter:

```
shared Absent|Value max<Value, Absent>(Iterable<Value, Absent> values)
  given Value satisfies Comparable<Value>
  given Absent satisfies Null { ... }
```

Exactly the right type pops out automatically. (And may be inferred.)

```
Null maxOfNone = max {}; //known to be empty
String maxOfSome = max { "hello", "world" }; //known to be nonempty

{String*} noneOrElse = ... ;
String? max = max(noneOrElse); //might be empty or nonempty
```

# Idiom #7

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## Idiom: multiple return values

For example, an operation might return a `Protocol` *and* a `Path`.

```
//Java
class ProtocolAndPath { ... }

ProtocolAndPath parseUrl(String url) {
    return new ProtocolAndPath(protocol(url), path(url));
}
```

We *have to* define a class.

# Idiom #7

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## Idiom: multiple return values

A function can be defined to return a tuple type.

```
[Protocol, Path] parseUrl(String url)
=> [protocol(url), path(url)];
```

Now a caller may extract the individual return values:

```
value protocolAndPath = parseUrl(url);
Protocol name = protocolAndPath[0];
Path address = protocolAndPath[1];
```

What about other indexes?

```
Null missing = protocolAndPath[3];
Protocol|Path|Null val = nameAndAddress[index];
```

# Idiom #8

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## Idiom: spreading tuple return values

Imagine we want to pass the result of `parseUrl()` to another function

```
Response get(Protocol name, Path address) => ... ;
```

We can use the spread operator, `*`, like in Groovy:

```
value response = get(*parseUrl(url));
```

Or we can work at the function level, using `unflatten()`

```
Response(String) get = compose(unflatten(get), parseUrl);  
value response = get("http://ceylon-lang.org");
```

There is a deep relationship between function types and tuple types.

# Idiom #9

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## Idiom: abstract over function types

Problem: the `compose()` function composes functions.

```
X(A) compose<X,Y,A>(X(Y) x, Y(A) y)
=> (A a) => x(y(a));
```

But this is not quite as general as it could be!

For functions with just one parameter it works well:

```
Anything(Float) printSqrt = compose(print,sqrt);
```

What about functions with multiple parameters?

```
value printSum = compose(print,plus);
```

# Idiom #9

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## Idiom: abstract over function types

Solution: abstract over unknown tuple type.

```
X(*Args) compose<X,Y,Args>(X(Y) x, Y(*Args) y)  
  given Args satisfies Anything[]  
  => flatten((Args args) => x(y(*args))));
```

A little uglier, but does the job!

```
Anything(Float,Float) printSum = compose(print,plus);
```

Even if this doesn't seem that useful at first sight, we actually use it in all sorts of places: for example, in the metamodel, and in `ceylon.promise`